# APEC 8001: Problem Set 1

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1. Proposition 1.B.1 in MWG states the following:

If  $\gtrsim$  is rational then:

- i. > is both irreflexive (x > x cannot hold) and transitive (if x > y and y > z, then x > z)
- ii. ~ is reflexive ( $x \sim x$  for all x), transitive (if  $x \sim y$  and  $y \sim z$ , then  $x \sim z$ ) and symmetric (if  $x \sim y$  then  $y \sim x$ ),
- iii. If  $x \succ y$  and  $y \gtrsim z$ , then  $x \succ z$

Provide a proof of the three results in this proposition. Hint: Showing that  $\sim$  is transitive is the easiest, so start there.

### Answer:

There are several ways to show these results, the proofs below are just one example:

To show  $\succ$  is irreflexive, using the definition of a strict preference relation  $\succ$ ,  $x \succ x \Leftrightarrow x \gtrsim x$  but not  $x \gtrsim x$ , yielding a contradiction. Thus,  $\succ$  is irreflexive.

To show > is transitive, using the definition of a strict preference relation >,  $\forall x, y, z \in X$  we have  $x > y \Leftrightarrow x \gtrsim y$  but not  $y \gtrsim x$  and  $y > z \Leftrightarrow y \gtrsim z$  but not  $z \gtrsim y$ . By transitivity of  $\gtrsim, x \gtrsim y$  and  $y \gtrsim z \Rightarrow x \gtrsim z$ . For purposes of contradiction, suppose  $z \gtrsim x$ . Then by transitivity of  $\gtrsim, z \gtrsim x$  and  $x \gtrsim y \Rightarrow z \gtrsim y$ . Since we cannot have  $z \gtrsim y$  if y > z, this yields a contradiction. So, it cannot be that  $z \gtrsim x$ . And  $x \gtrsim z$  but not  $z \gtrsim x \Leftrightarrow x > z$ . Thus, x > y and  $y > z \Rightarrow x > z$ , proving > is transitive.

To show ~ is reflexive, using the definition of an indifference relation ~,  $\forall x, y \in X$  we have  $x \sim y \Leftrightarrow x \gtrsim y$  and  $y \gtrsim x$ . By transitivity of  $\gtrsim$ ,  $x \gtrsim y$  and  $y \gtrsim x \Rightarrow x \gtrsim x$ . We know  $x \gtrsim x$  and  $x \gtrsim x \Rightarrow x \sim x$ . Thus, ~ is reflexive.

To show ~ is transitive, using the definition of an indifference relation ~,  $\forall x, y, z \in X$  we have  $x \sim y \Leftrightarrow x \gtrsim y$  and  $y \gtrsim x$  and  $y \sim z \Leftrightarrow y \gtrsim z$  and  $z \gtrsim y$ . By transitivity of  $\gtrsim, x \gtrsim y$  and  $y \gtrsim z \Rightarrow x \gtrsim z$ . Similarly,  $y \gtrsim x$  and  $z \gtrsim y \Rightarrow z \gtrsim x$ . We know  $x \gtrsim z$  and  $z \gtrsim x \Rightarrow x \sim z$ . Thus,  $x \sim y$  and  $y \sim z \Rightarrow x \sim z$ , proving ~ is transitive.

To show ~ is symmetric, using the definition of an indifference relation ~,  $\forall x, y \in X$  we have  $x \sim y \Leftrightarrow x \gtrsim y$  and  $y \gtrsim x$ . Thus, we must have  $y \gtrsim x$  and  $x \gtrsim y$  which implies  $y \sim x$ . So, we have  $x \sim y \Rightarrow y \sim x$ . Thus, ~ is symmetric.

To show x > y and  $y \ge z$  implies x > z, we will use the definition of strict preference relation >. We have  $x > y \Leftrightarrow x \ge y$  but not  $y \ge x$ . By transitivity of  $\ge$ ,  $x \ge y$  and  $y \ge z \Rightarrow x \ge z$ . For purposes of contradiction, suppose  $z \ge x$ . Then by transitivity of  $\ge$ ,  $z \ge x$  and  $y \ge z \Rightarrow y \ge x$ , yielding a

contradiction. So, it cannot be the case that  $z \gtrsim x$ . We know  $x \gtrsim z$  but not  $z \gtrsim x \Leftrightarrow x > z$ . Thus, x > y and  $y \gtrsim z \Rightarrow x > z$ .

- 2. Given the choice set  $X = \{x, y, z\}$  and the choice structure ( $\mathscr{B}, \mathcal{C}()$ ), where  $\mathscr{B} = \{\{x, y\}, \{y, z\}, \{x, z\}, \{x, y, z\}\}, \mathcal{C}(\{x, y\}) = \{x\}, \mathcal{C}(\{y, z\}) = \{z\}, \mathcal{C}(\{x, z\}) = \{x, z\}, and \mathcal{C}(\{x, y, z\}) = \{x, z\}$ :
  - a. Demonstrate this choice structure satisfies the weak axiom of revealed preferences (WARP)

## Answer:

Remember,  $(\mathcal{B}, \mathcal{C}())$  satisfies WARP if for some  $B \in \mathcal{B}$ , with  $x, y \in B$ , we have  $x \in \mathcal{C}(B)$ , then for any  $B' \in \mathcal{B}$ , with  $x, y \in B'$  and  $y \in \mathcal{C}(B')$  we must also have  $x \in \mathcal{C}(B')$ . Let us consider each choice rule above:

 $C({x, y}) = {x}$ : Any choice rule over a different budget set containing both x and y must include x (if it includes y)

 $C(\{y, z\}) = \{z\}$ : Any choice rule over a different budget set containing both y and z must include z (if it includes y)

 $C({x, z}) = {x, z}$ : Any choice rule over a different budget set containing both x and z must include both x and z

 $C(\{x, y, z\}) = \{x, z\}$ : Any choice rule over a different budget set containing x, y and z must include both x and z (if it includes y). Also, any choice rule over a different budget set containing both x and y must include x (if it includes y), any choice rule over a different budget set containing both y and z must include z (if it includes y, and any choice rule over a different budget set containing both x and z must include z (if it includes y, and any choice rule over a different budget set containing both x and z must include both x and z.

All the choice rules in this choice structure are consistent, and thus WARP is satisfied.

b. Suppose instead  $C(\{x, y, z\}) = \{x\}$ . Does this choice structure satisfy WARP? Provide an explanation for your answer.

### Answer:

Again, let us consider each choice rule in this choice structure:

 $C({x, y}) = {x}$ : Any choice rule over a different budget set containing both x and y must include x (if it includes y)

 $C(\{y, z\}) = \{z\}$ : Any choice rule over a different budget set containing both y and z must include z (if it includes y)

 $C({x, z}) = {x, z}$ : Any choice rule over a different budget set containing both x and z must include both x and z

 $C({x, y, z}) = {x}$ : Any choice rule over a different budget set containing x, y and z must include x (if it includes y or z). Also, any choice rule over a different budget set containing both x and y must include x

(if it includes y), and any choice rule over a different budget set containing both x and z must include both x (if it contains z).

Because  $C(\{x, y, z\})$  only contains x, when  $C(\{x, z\})$  contains both x and z, WARP is violated.

Consider a consumer's choice of spending wealth on only two goods, x1 and x2. Show, for each of the following Walrasian demands, that they satisfy: i) homogeneity of degree zero, and ii) Walras' law:

a. 
$$x_1(p_1, p_2, w) = \frac{w}{3p_1}$$
 and  $x_2(p_1, p_2, w) = \frac{2w}{3p_2}$ 

Answer:

$$x_1(\tau p_1, \tau p_2, \tau w) = \frac{\tau w}{3\tau p_1} = \frac{w}{3p_1} = x_1(p_1, p_2, w)$$
$$x_2(\tau p_1, \tau p_2, \tau w) = \frac{2\tau w}{3\tau p_2} = \frac{2w}{3p_2} = x_2(p_1, p_2, w)$$

Thus, the Walrasian demands are homogenous of degree zero

$$p_1 x_1(p_1, p_2, w) + p_2 x_2(p_1, p_2, w) = p_1 \left(\frac{w}{3p_1}\right) + p_2 \left(\frac{2w}{3p_2}\right) = \frac{w}{3} + \frac{2w}{3} = w$$

Thus, the Walrasian demands satisfy Walras' law

b. 
$$x_1(p_1, p_2, w) = \frac{p_2^2}{4p_1^2}$$
 and  $x_2(p_1, p_2, w) = \frac{w}{p_2} - \frac{p_2}{4p_1}$ 

#### Answer:

$$x_{1}(\tau p_{1}, \tau p_{2}, \tau w) = \frac{(\tau p_{2})^{2}}{4(\tau p_{1})^{2}} = \frac{\tau^{2} p_{2}^{2}}{4\tau^{2} p_{1}^{2}} = \frac{p_{2}^{2}}{4p_{1}^{2}} = x_{1}(p_{1}, p_{2}, w)$$
$$x_{2}(\tau p_{1}, \tau p_{2}, \tau w) = \frac{\tau w}{\tau p_{2}} - \frac{\tau p_{2}}{4\tau p_{1}} = \frac{w}{p_{2}} - \frac{p_{2}}{4p_{1}} = x_{2}(p_{1}, p_{2}, w)$$

Thus, the Walrasian demands are homogenous of degree zero

$$p_1 x_1(p_1, p_2, w) + p_2 x_2(p_1, p_2, w) = p_1 \left(\frac{p_2^2}{4p_1^2}\right) + p_2 \left(\frac{w}{p_2} - \frac{p_2}{4p_1}\right) = \frac{p_2^2}{4p_1} + w - \frac{p_2^2}{4p_1} = w$$

Thus, the Walrasian demands satisfy Walras' law

c. 
$$x_1(p_1, p_2, w) = \frac{w}{p_1 + p_2}$$
 and  $x_2(p_1, p_2, w) = \frac{w}{p_1 + p_2}$ 

#### Answer:

$$x_1(\tau p_1, \tau p_2, \tau w) = \frac{\tau w}{\tau p_1 + \tau p_2} = \frac{\tau w}{\tau (p_1 + p_2)} = \frac{w}{p_1 + p_2} = x_1(p_1, p_2, w)$$
$$x_2(\tau p_1, \tau p_2, \tau w) = \frac{\tau w}{\tau p_1 + \tau p_2} = \frac{\tau w}{\tau (p_1 + p_2)} = \frac{w}{p_1 + p_2} = x_2(p_1, p_2, w)$$

Thus, the Walrasian demands are homogenous of degree zero

$$p_1 x_1(p_1, p_2, w) + p_2 x_2(p_1, p_2, w) = p_1 \left(\frac{w}{p_1 + p_2}\right) + p_2 \left(\frac{w}{p_1 + p_2}\right) = \frac{p_1 w}{p_1 + p_2} + \frac{p_2 w}{p_1 + p_2} = \frac{w(p_1 + p_2)}{p_1 + p_2} = w$$

Thus, the Walrasian demands satisfy Walras' law